

Bipath Persistent homology and its stability

Shunsuke Tada

Mathematical Science Center for Co-creative Society

Tohoku University

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The aim is to introduce

Bipath persistent homology

which is an extension of persistent homology with a visualization (bipath persistence diagram) and

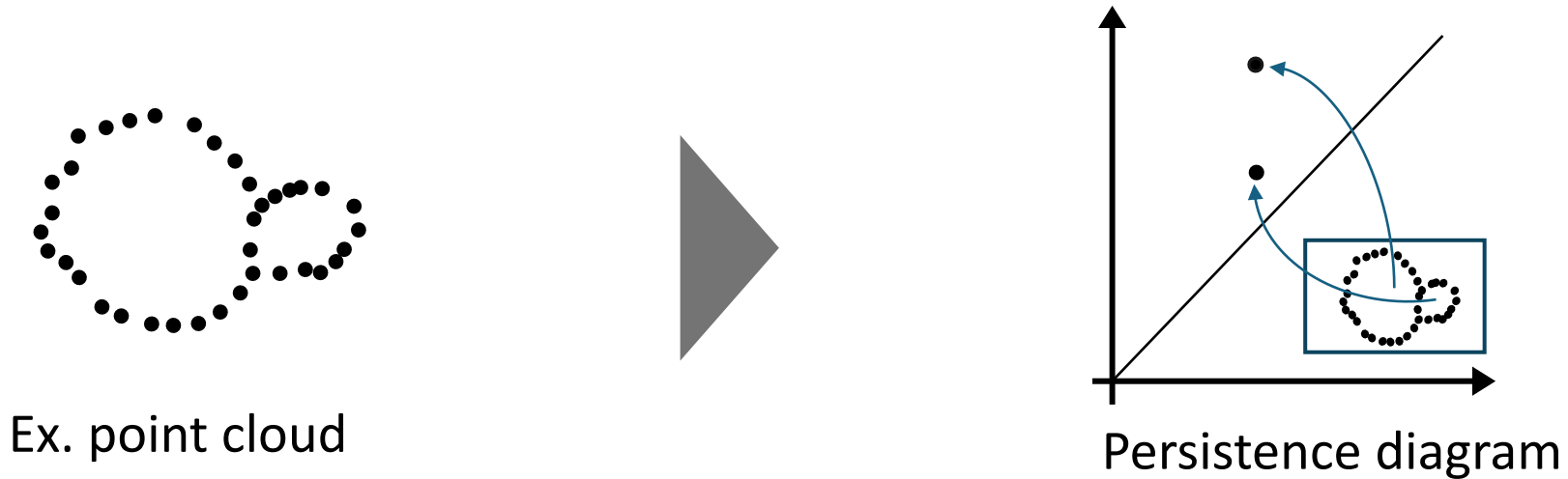
stability properties

Contents

- (1) Introduction: Persistent homology and related settings.
- (2) Stability property of bipath persistence: Isometry theorem.

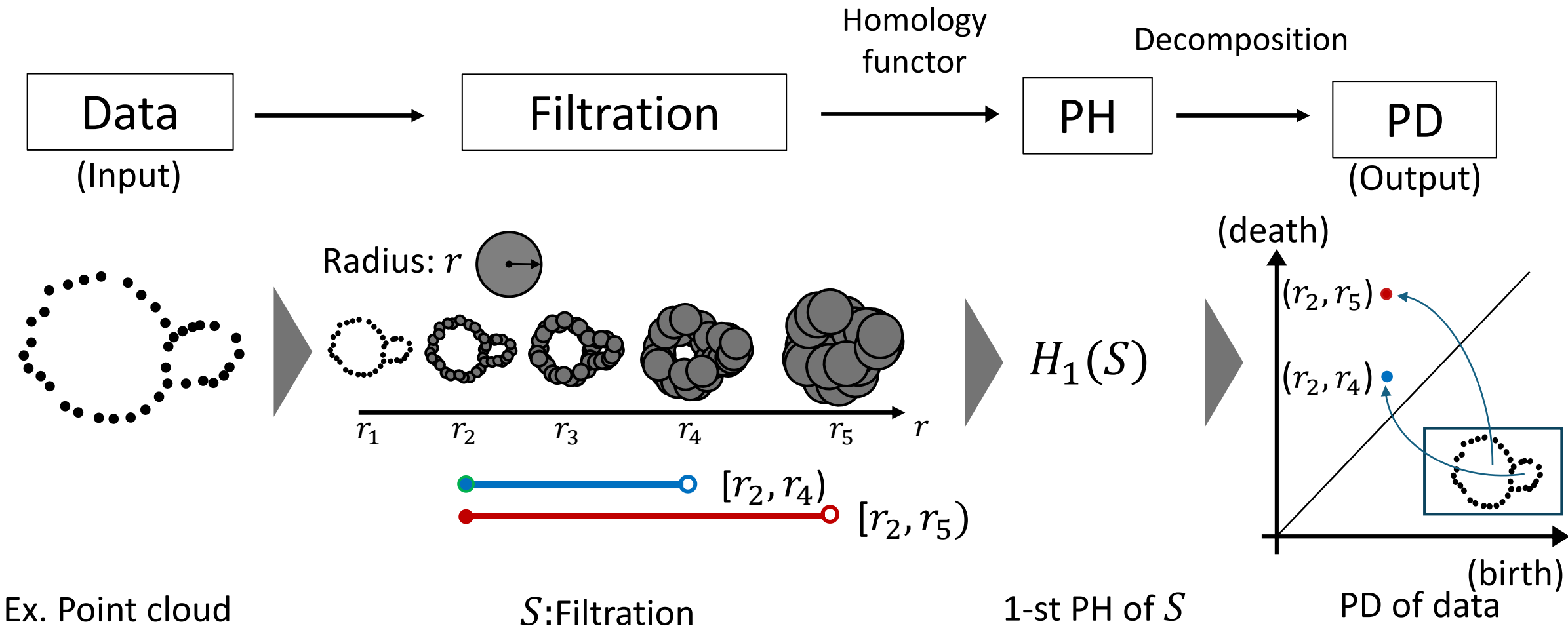
Introduction

- *Persistent homology* (PH) is a tool in Topological Data Analysis.
- It captures the persistence of “shape” (connected components, holes or voids) of data by a *persistence diagram* (PD).

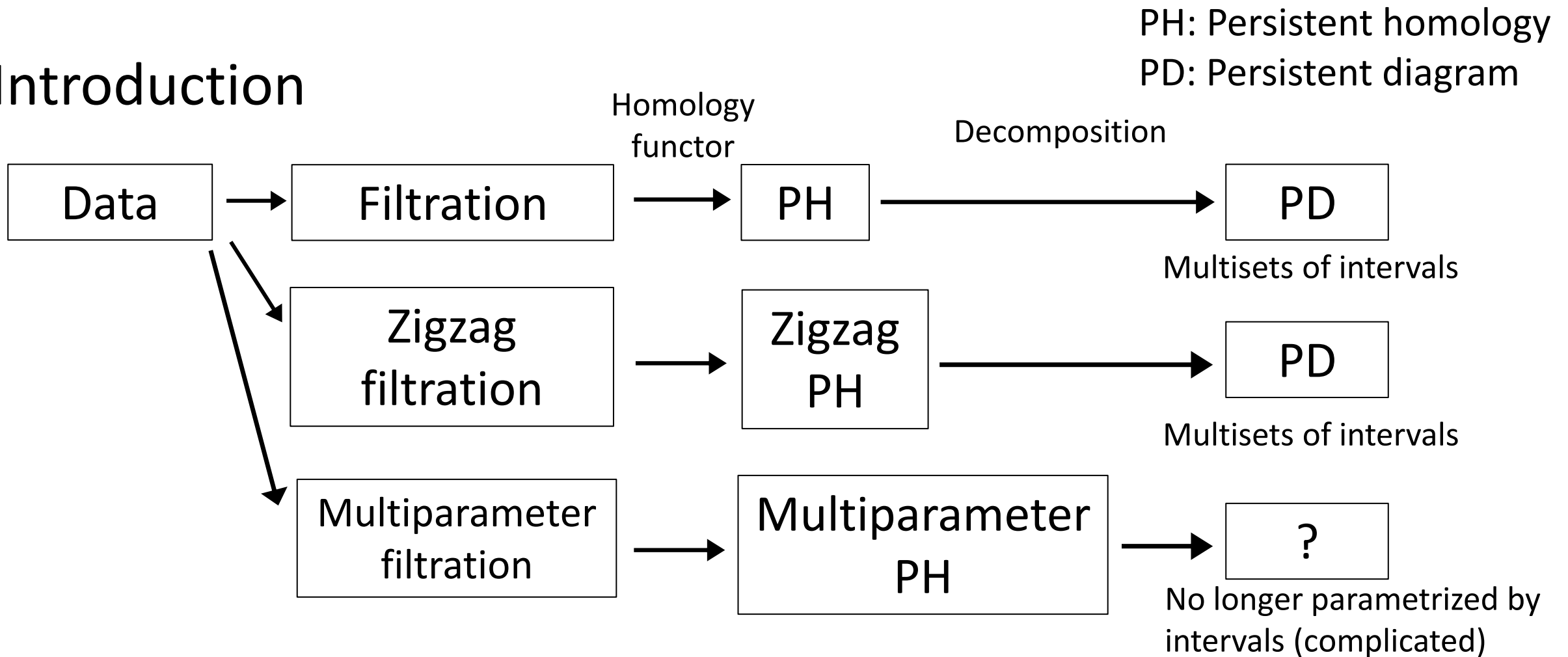


Introduction

PH: Persistent homology
PD: Persistent diagram

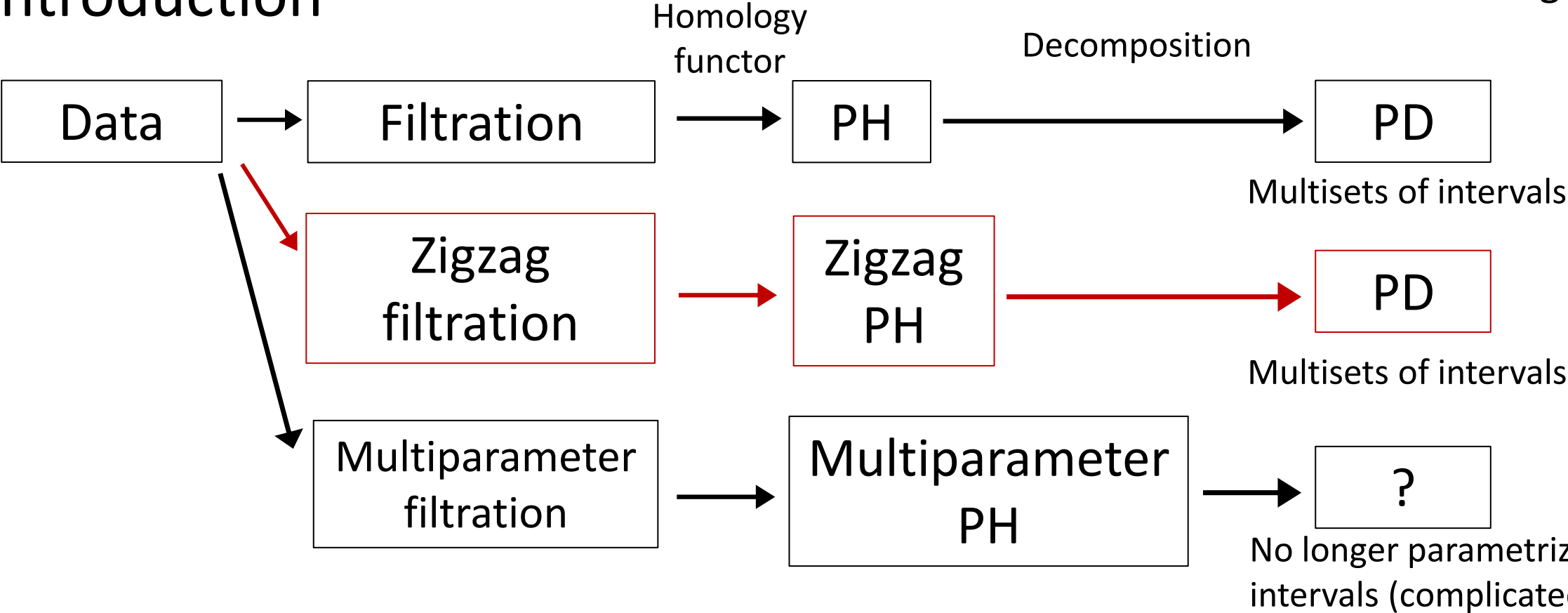


Introduction



Introduction

PH: Persistent homology
 PD: Persistent diagram



$$S: S_1 \supset S_2 \subset S_3 \supset S_4 \subset S_5$$

Zigzag filtration

➤

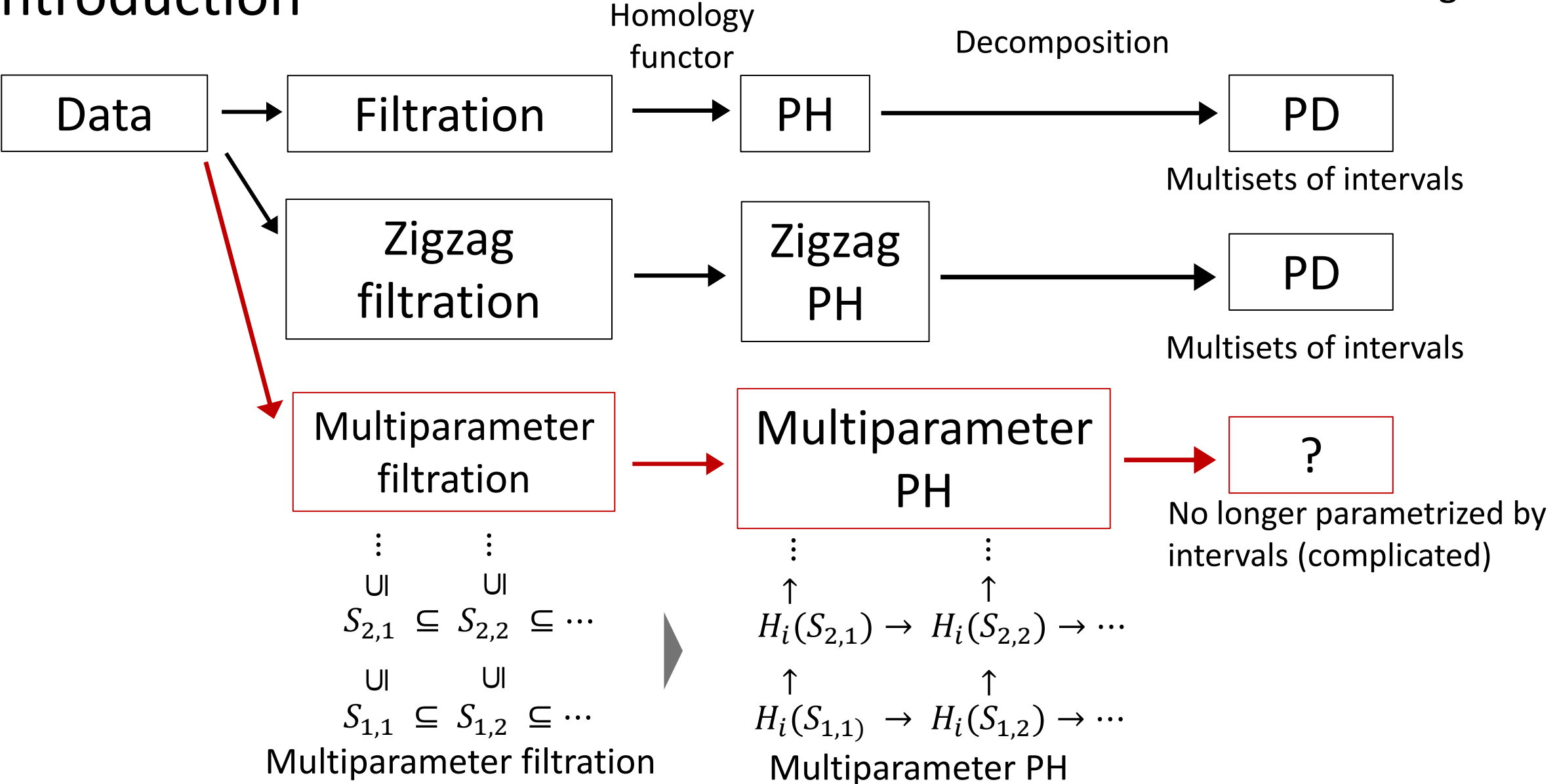
$$H_i(S) = H_i(S_{r_1}) \leftarrow H_i(S_{r_2}) \rightarrow H_i(S_{r_3}) \leftarrow H_i(S_{r_4}) \rightarrow H_i(S_{r_5})$$

(Interval-decomposable)

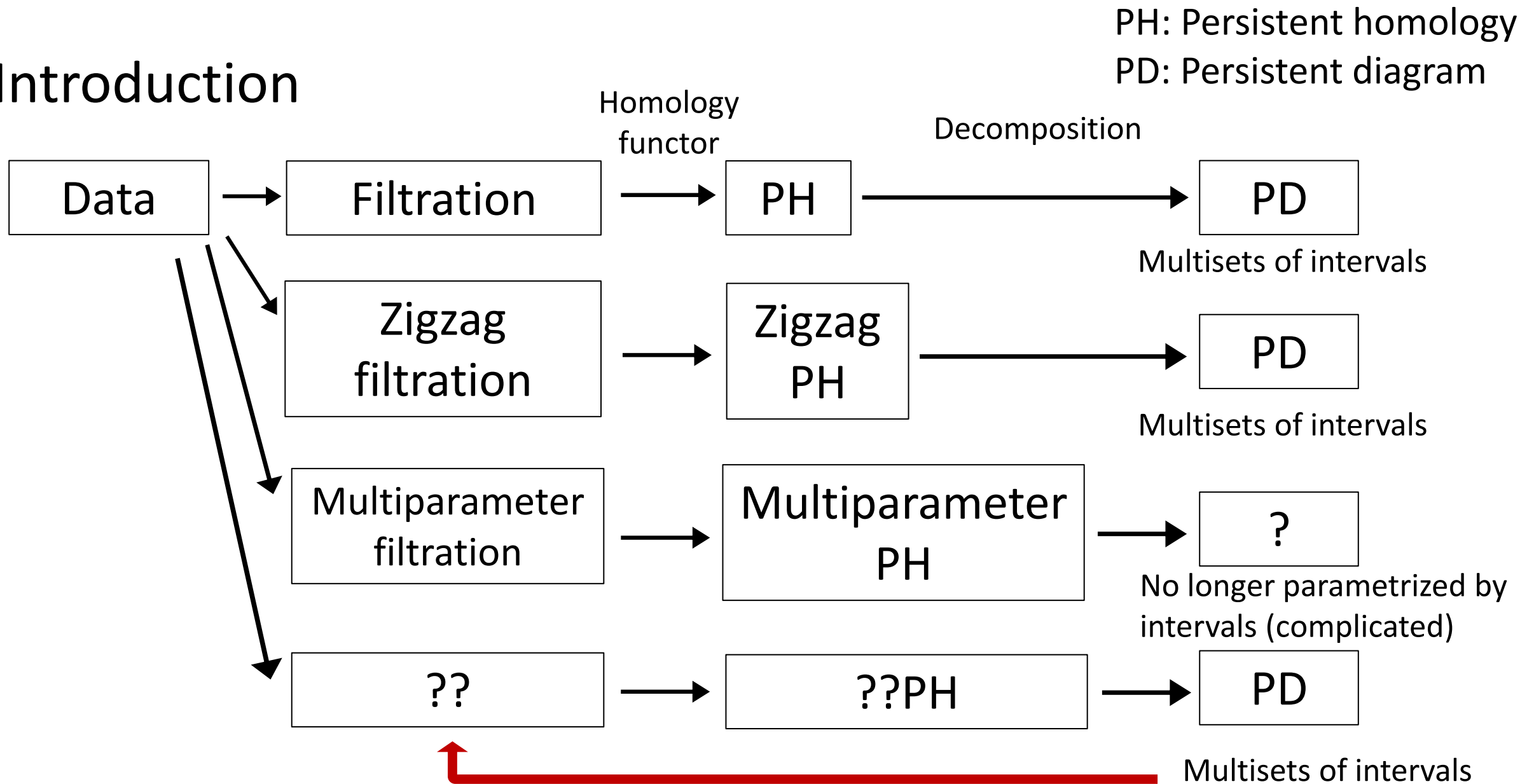
Zigzag PH

Introduction

PH: Persistent homology
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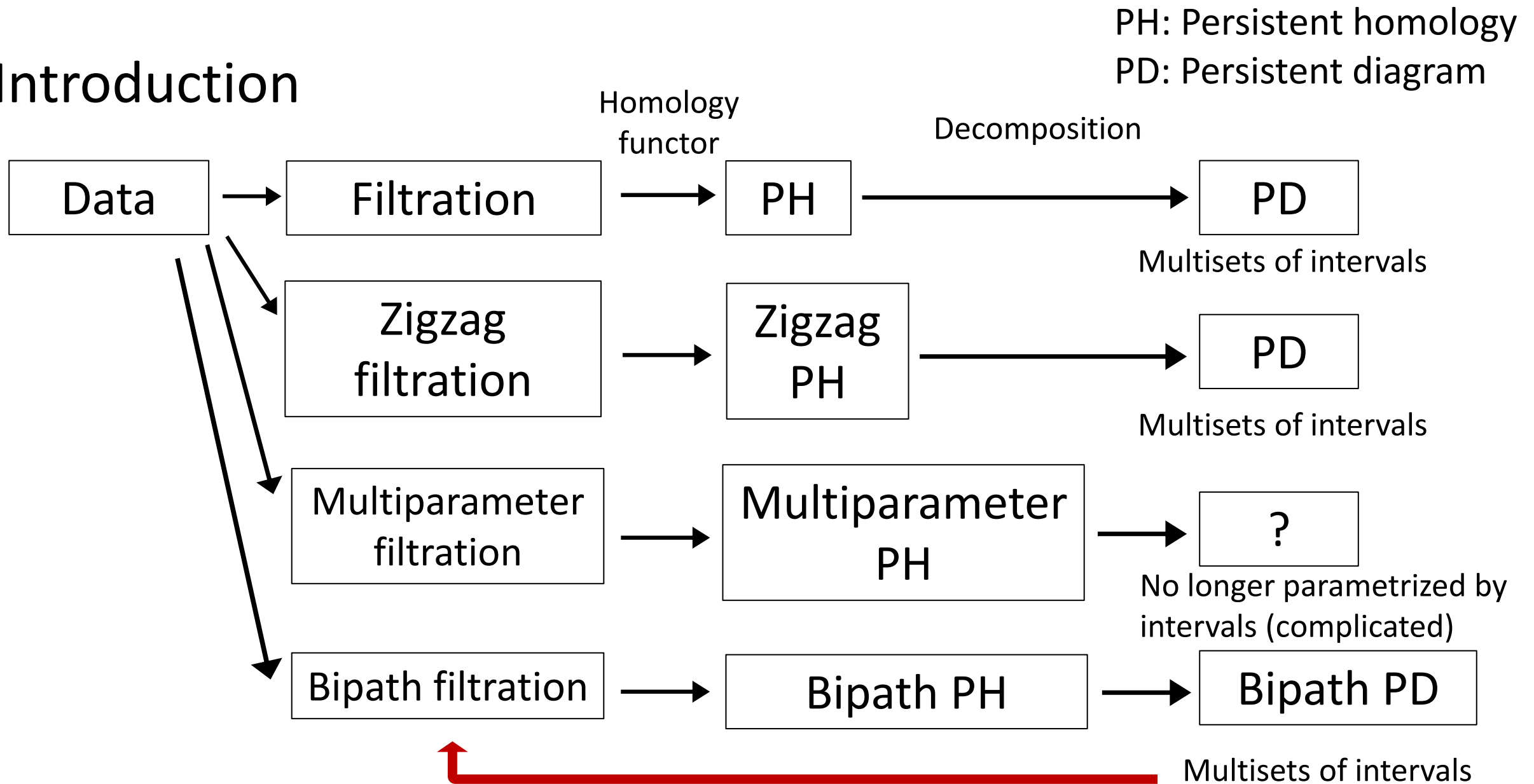


Introduction



Do we have other arrangement of spaces like standard/zigzag filtration?

Introduction



We propose bipath persistent homology as a new framework.

Introduction

Theorem [Aoki-Escolar-T, 25]

Let P be a connected finite poset. The following are equivalent.

(a) Every P -persistence module V is interval-decomposable.

(b) The Hasse diagram of P is one of the following forms:

$$A_n(a): 1 \longleftrightarrow \cdots \longleftrightarrow n$$

Zigzag posets (type A)

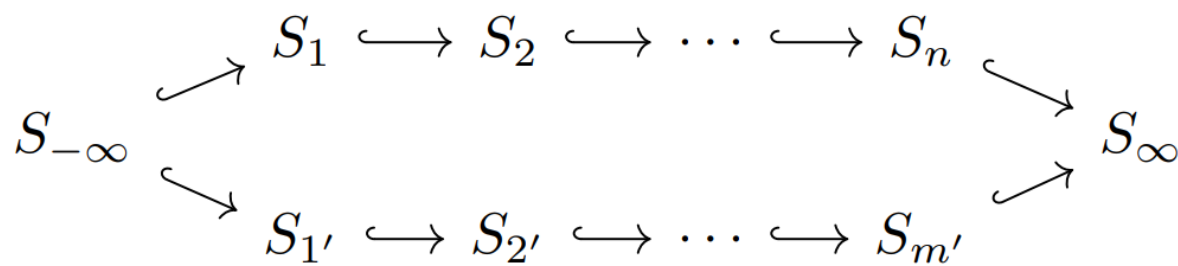
$$B_{n,m}: \begin{array}{c} \nearrow 1 \rightarrow 2 \xrightarrow{\leq} \cdots \rightarrow n \searrow \\ -\infty \qquad \qquad \qquad +\infty \\ \searrow \qquad \qquad \qquad \nearrow \\ 1' \rightarrow 2' \rightarrow \cdots \rightarrow m' \end{array}$$

Bipath posets

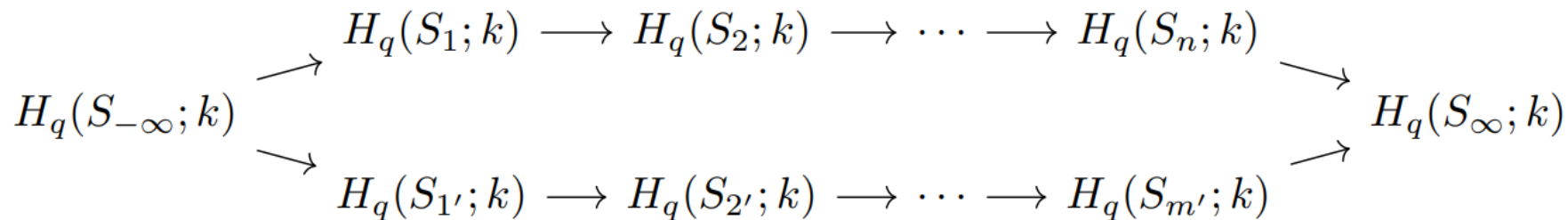
$x \leftrightarrow y$ represents either $x \rightarrow y$ or $x \leftarrow y$.

Introduction

We can consider a *bipath persistent homology* (bipath PH) of a *bipath filtration*, which can capture topological features across the filtration.



Bipath filtration

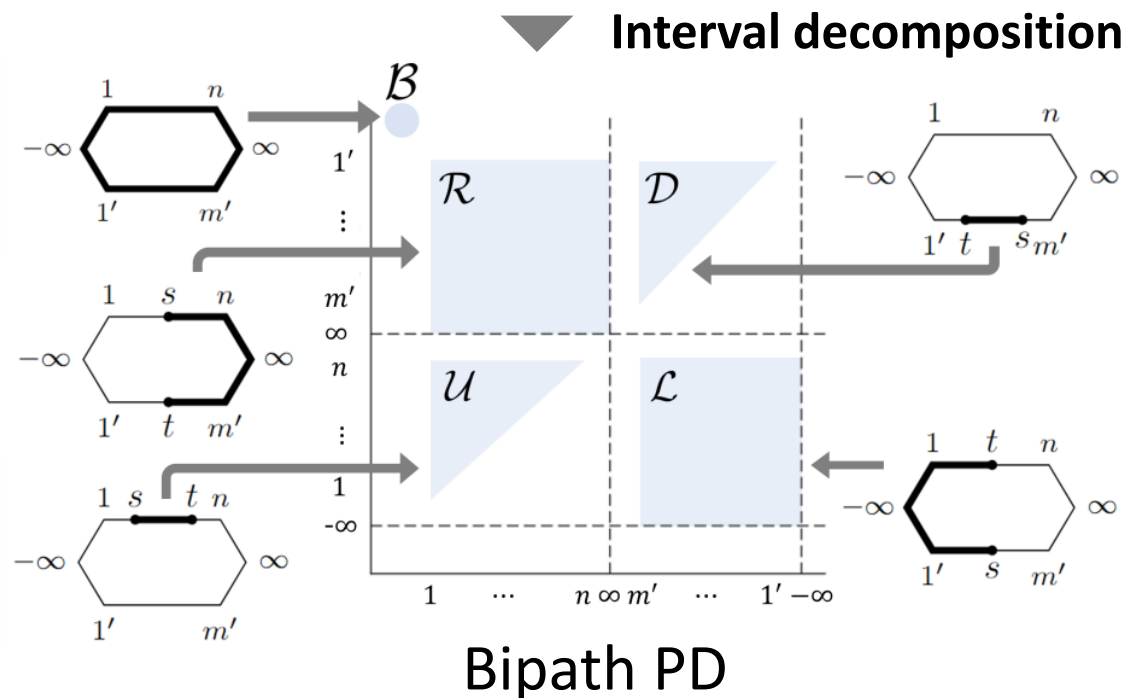


Bipath PH

Introduction

We can get a *Bipath Persistence Diagram* (Bipath PD).

$$\begin{array}{c}
 H_q(S_{-\infty}; k) \begin{array}{l} \nearrow H_q(S_1; k) \longrightarrow H_q(S_2; k) \longrightarrow \cdots \longrightarrow H_q(S_n; k) \\ \searrow H_q(S_{1'}; k) \longrightarrow H_q(S_{2'}; k) \longrightarrow \cdots \longrightarrow H_q(S_{m'}; k) \end{array} \begin{array}{l} \nearrow \\ \searrow \end{array} H_q(S_{\infty}; k)
 \end{array}$$



T. Aoki, E. G. Escobar, and S. Tada.
Bipath persistence. Japan Journal of Industrial
and Applied Mathematics, 42:453–486, 2025.

Introduction

A recent study on bipath persistent homology:

	Bipath PH
Interval-Decomposability	○
Visualization(Bipath PD)	○
Algorithm(Implementation)	○
Stability theorem for Bipath PD	○
Application	-

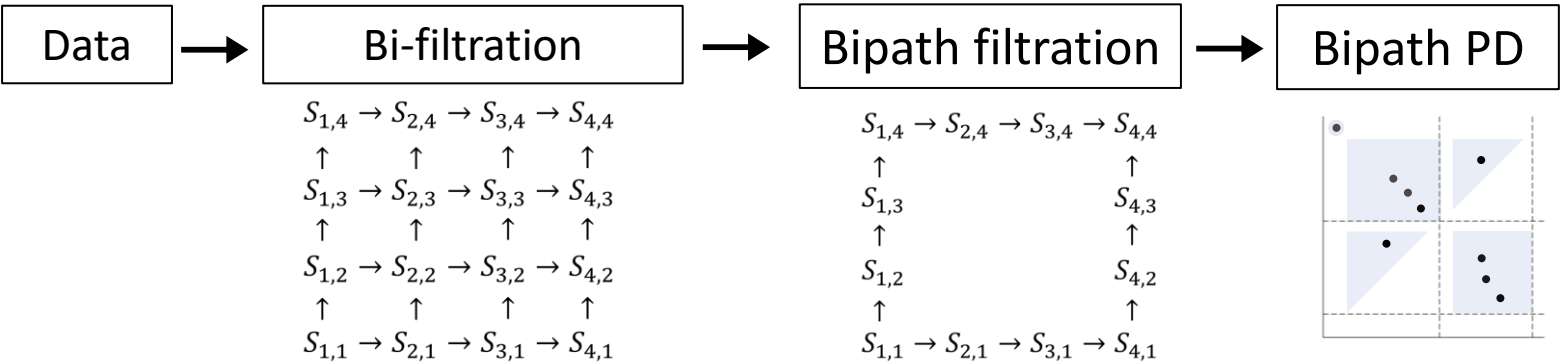
Algorithm: Aoki, T., Escolar, E.G. & Tada, S. Bipath persistence. *Japan J. Indust. Appl. Math.* **42**, 453–486 (2025).

Implementation: <https://github.com/ShunsukeTada1357/Bipathposets>

Introduction

A recent study on bipath persistent homology:

	Bipath PH
Interval-Decomposability	○
Visualization(Bipath PD)	○
Algorithm(Implementation)	○
Stability theorem for Bipath PD	○
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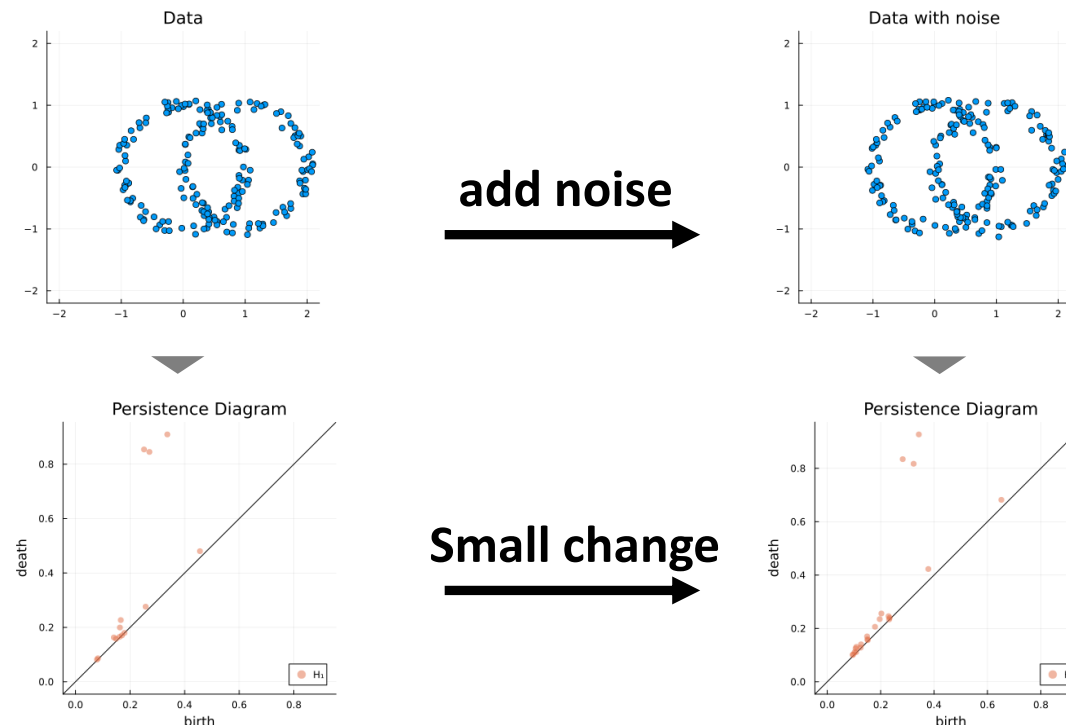
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Stability of bipath persistence diagrams

Background: Stability theorem for standard PH

- In persistent homology analysis stability theorem [Cohen-Steiner–Edelsbrunner–Harer '07] is important.
- This implies small changes in data implies small changes in the PD.
 \leadsto It justifies the use of PH for studying noisy data.

Example



Background: Stability theorem for standard PH

Recall that stability theorem can be deduced by the isometry theorem.

Isometry theorem [Lesnick '15]

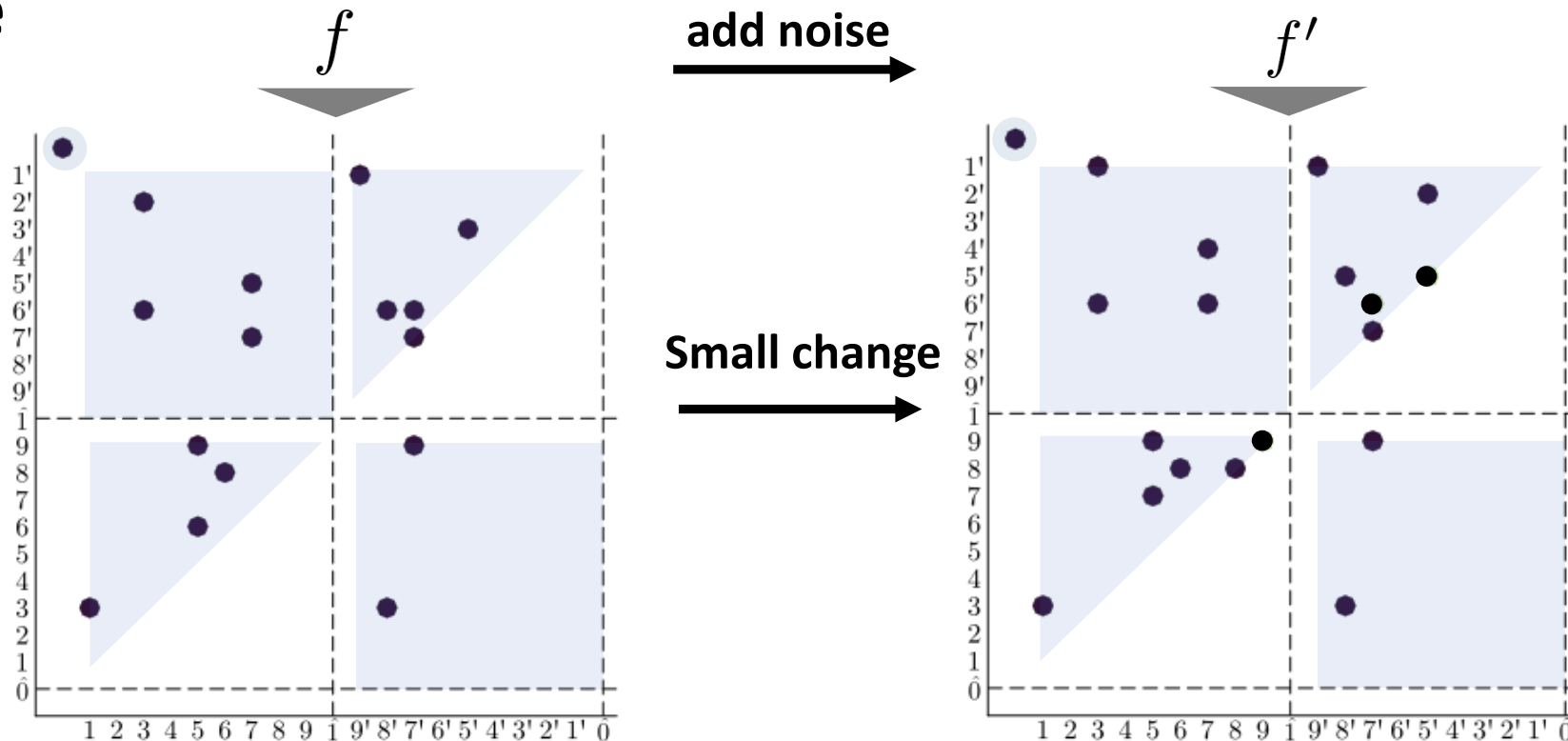
Let V and W be \mathbb{R} -persistence modules. Then, V and W are ϵ -interleaved if and only if there exist an ϵ -matching between $\mathcal{B}(V)$ and $\mathcal{B}(W)$. Thus, we have

$$d_B(\mathcal{B}(V), \mathcal{B}(W)) = d_I(V, W).$$

Stability theorem for bipath

Stability of bipath PDs holds [T '25, Theorem 4.1].

Example



- This suggests a justification for using bipath PH on noisy data.
- This can be deduced by the isometry theorem for bipath persistence.

Stability theorem for bipath

To discuss stability, we consider continuous bipath poset B .

$$B_{n,m}: \begin{array}{c} 1 \rightarrow 2 \rightarrow \dots \rightarrow n \\ \nearrow \quad \searrow \\ -\infty \quad \quad +\infty \\ \searrow \quad \nearrow \\ 1' \rightarrow 2' \rightarrow \dots \rightarrow m' \end{array} \quad \leadsto \quad B: \begin{array}{c} \mathbb{R} \times \{1\} \\ \nearrow \quad \searrow \\ -\infty \quad \quad +\infty \\ \searrow \quad \nearrow \\ \mathbb{R} \times \{2\} \end{array} \\ (\mathbb{R} \times \{1\}) \sqcup (\mathbb{R} \times \{2\}) \sqcup \{-\infty, +\infty\}$$

Isometry theorem for bipath persistence [T '25]

Let V and W be B -persistence modules. Then V and W are ϵ -interleaved if and only if there exist an ϵ -matching between $\mathcal{B}(V)$ and $\mathcal{B}(W)$. Thus, we have $d_B(\mathcal{B}(V), \mathcal{B}(W)) = d_I(V, W)$.

- \leadsto
- Setting for the definitions of d_I and d_B .
 - Graph-theoretic approach in the general setting.
 - Return to the bipath setting.

Stability theorem: Setting

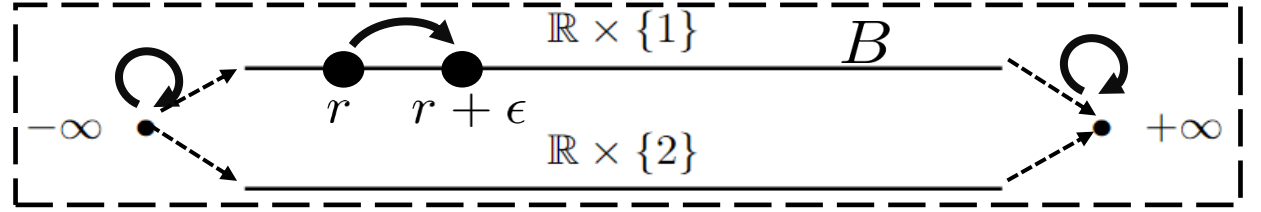
- Let k be a field, and let P be a poset.
- A P -persistence module is an object in $\text{rep}_k(P) := \text{Fun}(P, \text{vect}_k)$.
- For $V \cong \bigoplus_{\gamma \in \Gamma} V_\gamma \in \text{rep}_k(P)$ (V_γ : indecomposable), set $\mathcal{B}(V) := \{ V_\gamma \mid \gamma \in \Gamma \}$.
- A translation on P is an order-preserving map $h: P \rightarrow P$ s. t. $p \leq h(p)$ for every $p \in P$.
- Fix a family of translations $\Lambda := \{ \Lambda_\epsilon \}_{\epsilon \in \mathbb{R}_{\geq 0}}$ on P satisfying:

$$\Lambda_0 = \text{id}_P \text{ and } \Lambda_{\epsilon+\zeta} = \Lambda_\epsilon \circ \Lambda_\zeta \text{ for all } \epsilon, \zeta \in \mathbb{R}_{\geq 0}.$$

Example [T, '25, Definition 3.4].

Let B be the bipath poset. We define $\Lambda_\epsilon^B := \{ \Lambda_\epsilon^B \}_{\epsilon \in \mathbb{R}_{\geq 0}}$ by

$\Lambda_\epsilon^B(\pm\infty) := \pm\infty$, and $\Lambda_\epsilon^B((r, i)) := (r + \epsilon, i)$ for $(r, i) \in \mathbb{R} \times \{i\}$ ($i = 1, 2$).



Stability theorem: Setting

- Let k be a field, and let P be a poset.
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Interleaving and bottleneck distances are defined w. r. t. $\Lambda := \{ \Lambda_\epsilon \}_{\epsilon \in \mathbb{R}_{\geq 0}}$.

Stability theorem: Setting

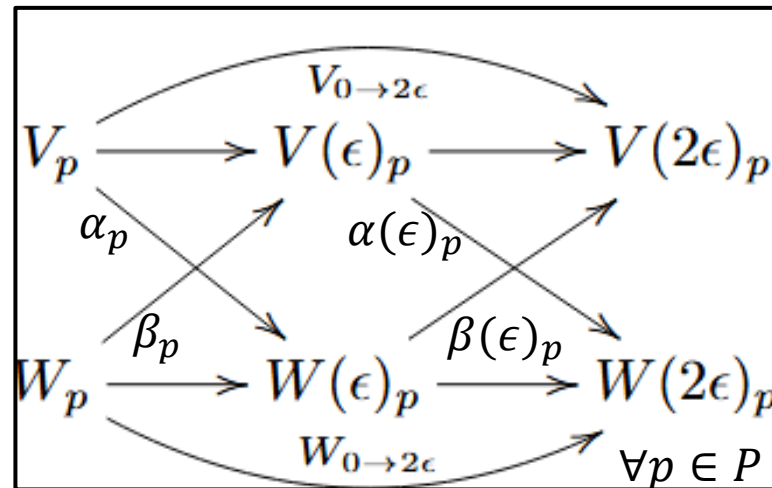
Let V, W be P -persistence modules, and $\epsilon \geq 0$.

- We write $V(\epsilon) := V \circ \Lambda_\epsilon \in \text{rep}_k(P)$

(this gives a functor $(\cdot)(\epsilon): \text{rep}_k(P) \rightarrow \text{rep}_k(P)$),
then, we have the induced morphism $V_{0 \rightarrow \epsilon}: V \rightarrow V(\epsilon)$.

- We say that V and W are ϵ -interleaved and write $V \sim_\epsilon W$ if there is a pair of morphisms $\alpha: V \rightarrow W(\epsilon)$ and $\beta: W \rightarrow V(\epsilon)$ s. t.

$$V_{0 \rightarrow 2\epsilon} = \beta(\epsilon) \circ \alpha \text{ and } W_{0 \rightarrow 2\epsilon} = \alpha(\epsilon) \circ \beta.$$



Stability theorem: Setting

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$$V_{0 \rightarrow 2\epsilon} = \beta(\epsilon) \circ \alpha \text{ and } W_{0 \rightarrow 2\epsilon} = \alpha(\epsilon) \circ \beta.$$

Definition (Interleaving distance)

The interleaving distance between P -persistence modules V and W is defined by $d_I^\Lambda(V, W) := \inf \{ \epsilon \in \mathbb{R}_{\geq 0} \mid V \sim_\epsilon W \}$.

Stability theorem: Setting

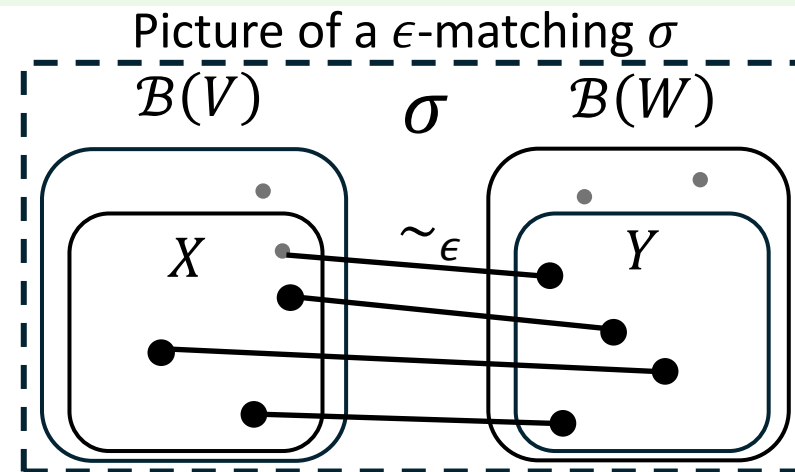
We say that a P -persistence module V is ϵ -trivial if $V_{0 \rightarrow \epsilon} = 0$.

Definition (ϵ -matching)

Let V and W be P -persistence modules. An ϵ -matching between $\mathcal{B}(V)$ and $\mathcal{B}(W)$ is a partial matching $\sigma: \mathcal{B}(V) \supseteq X \xrightarrow{1:1} Y \subseteq \mathcal{B}(W)$ satisfying:

- Every $I \in (\mathcal{B}(V) \sqcup \mathcal{B}(W)) \setminus (X \sqcup Y)$ is 2ϵ -trivial.
- If $\sigma(I) = J$, then $I \sim_{\epsilon} J$.

We say that V and W are ϵ -matched if there is an ϵ -matching.



- : 2ϵ -trivial
- : 2ϵ -**non**-trivial

Stability theorem: Setting

We say that a P -persistence module V is ϵ -trivial if $V_{0 \rightarrow \epsilon} = 0$.

Definition (ϵ -matching)

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- If $\sigma(I) = J$, then $I \sim_{\epsilon} J$.

We say that V and W are ϵ -matched if there is an ϵ -matching.

Definition (Bottleneck distance)

The bottleneck distance between P -persistence modules V and W is defined by $d_B^{\Lambda}(\mathcal{B}(V), \mathcal{B}(W)) := \inf \{ \epsilon \in \mathbb{R}_{\geq 0} \mid V \text{ and } W \text{ are } \epsilon\text{-matched} \}$.

Stability theorem: Outline.

Remark

Let V and W be P -persistence modules. If V and W are ϵ -matched, then they are ϵ -interleaved. Thus, we have

$$d_B^\Lambda(\mathcal{B}(V), \mathcal{B}(W)) \geq d_I^\Lambda(V, W).$$

(\because An ϵ -matching induces an ϵ -interleaving.)

\leadsto We observe the converse:

- V and W are ϵ -interleaved. $\Rightarrow V$ and W are ϵ -matched.

Step1: Interpreting an ϵ -matching as a matching in a bipartite graph

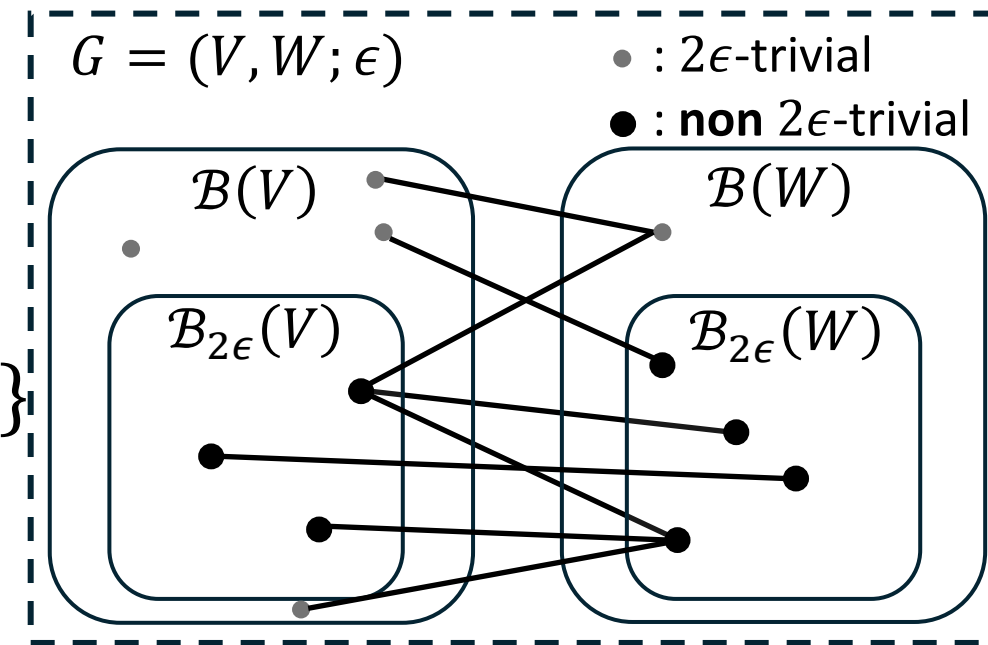
Step2: A sufficient condition for an ϵ -matching using a bipartite graph

Step3: Hall's marriage theorem is useful for showing the sufficient condition.

Step4: In the bipath setting, Step 2 is proved through Step 3.

Stability theorem: Outline Step 1

- Let V, W be P -persistence modules.
- Make bipartite graph $G = (V, W; \epsilon)$.
 - Vertices $\mathcal{B}(V) \sqcup \mathcal{B}(W)$
 - Edges $\{\{I, J\} \mid I \in \mathcal{B}(V), J \in \mathcal{B}(W), I \sim_{\epsilon} J\}$
- $\mathcal{B}_{2\epsilon}(V) := \{I \in \mathcal{B}(V) \mid I \text{ is **not** } 2\epsilon\text{-trivial}\}$



Proposition (1)[Bjerkevik '21, p.4]

The following are equivalent.

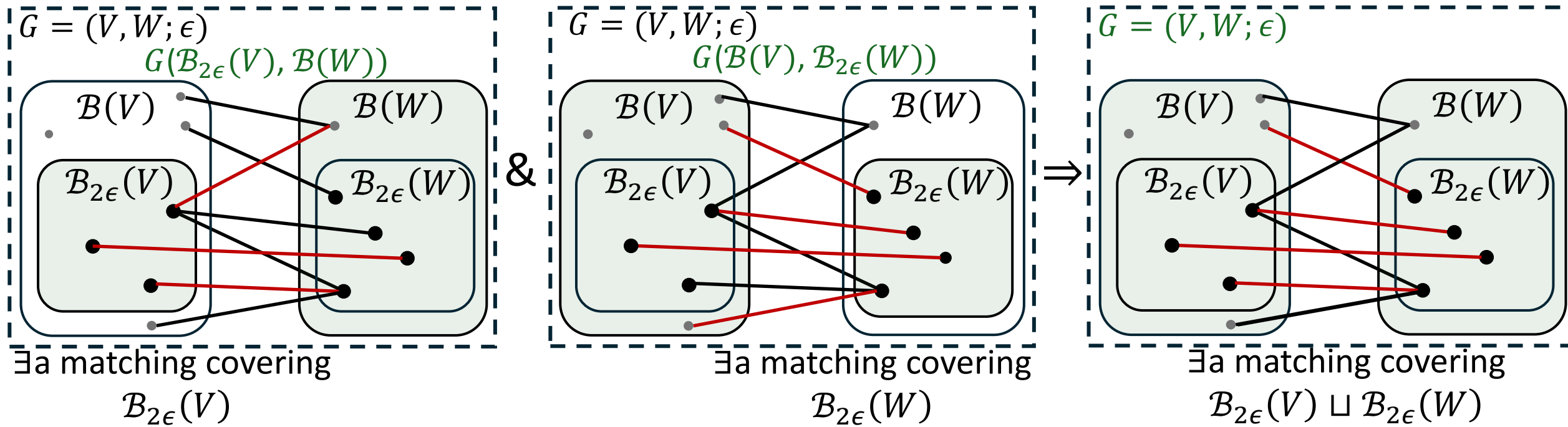
- V and W are ϵ -matched.
- \exists a matching in $G = (V, W; \epsilon)$ that covers $\mathcal{B}_{2\epsilon}(V) \sqcup \mathcal{B}_{2\epsilon}(W)$.

Stability theorem: Outline Step 2

Proposition (2) [cf. Bjerkevik, '21, p. 111]

Let V and W be P -persistence modules. If the following are satisfied, then there is a matching in $G = (V, W; \epsilon)$ that covers $\mathcal{B}_{2\epsilon}(V) \sqcup \mathcal{B}_{2\epsilon}(W)$

- \exists a matching in the full subgraph $G(\mathcal{B}_{2\epsilon}(V), \mathcal{B}(W))$ that covers $\mathcal{B}_{2\epsilon}(V)$.
- \exists a matching in the full subgraph $G(\mathcal{B}(V), \mathcal{B}_{2\epsilon}(W))$ that covers $\mathcal{B}_{2\epsilon}(W)$.

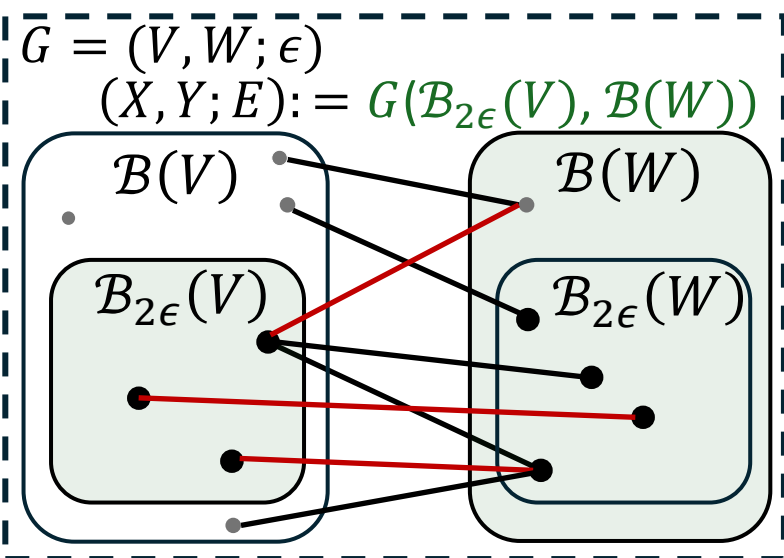


Stability theorem: Outline Step 3

Theorem (3) [Hall, 1935, Theorem 1]

Let $G = (X, Y; E)$ be a bipartite graph such that each vertex $x \in X$ has a finite neighborhood $N_G(x) \subseteq Y$. Then the following are equivalent:

- (a) \exists a matching in G that covers X .
- (b) For every finite subset $X' \subseteq X$, we have $|X'| \leq |\cup_{x \in X'} N_G(x)|$.



\leftarrow Since $V \in \text{rep}_k(P)$ is pointwise finite dimensional, $N_G(x) < \infty$ for every $x \in \mathcal{B}_{2\epsilon}(V)$ [Bje21, p.110].

$\leadsto (X, Y; E) := G(\mathcal{B}_{2\epsilon}(V), \mathcal{B}(W))$ satisfies the assumption of Hall's theorem.

\leadsto Existence of a matching covering $\mathcal{B}_{2\epsilon}(V)$ is equivalent to (b):

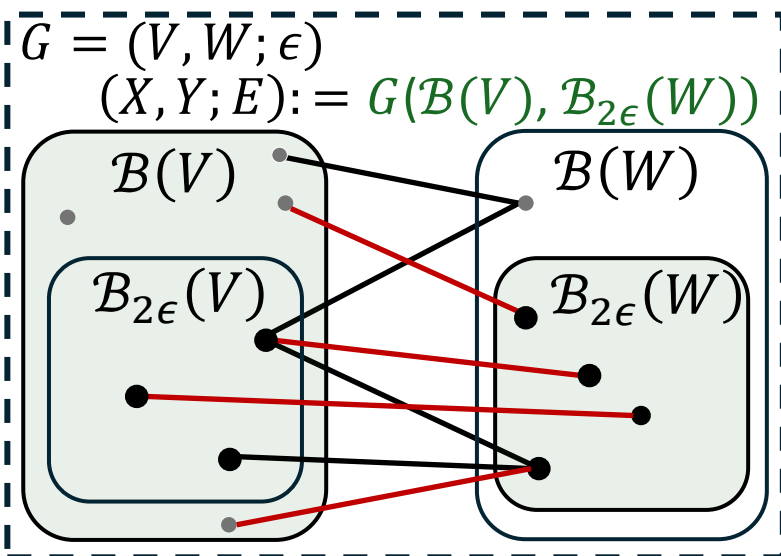
$$\forall X' \subseteq_{\text{fin.}} \mathcal{B}_{2\epsilon}(V), \text{ we have } |X'| \leq |\cup_{x \in X'} N_G(x)|.$$

Stability theorem: Outline Step 3

Theorem (3) [Hall, 1935, Theorem 1]

Let $G = (X, Y; E)$ be a bipartite graph such that each vertex $x \in X$ has a finite neighborhood $N_G(x) \subseteq Y$. Then the following are equivalent:

- (a) \exists a matching in G that covers X .
- (b) For every finite subset $X' \subseteq X$, we have $|X'| \leq |\cup_{x \in X'} N_G(x)|$.



\leftarrow Since $W \in \text{rep}_k(P)$ is pointwise finite dimensional, $N_G(x) < \infty$ for every $x \in \mathcal{B}_{2\epsilon}(W)$ [Bje21, p.110].

$\leadsto (X, Y; E) := G(\mathcal{B}(V), \mathcal{B}_{2\epsilon}(W))$ satisfies the assumption of Hall's theorem.

\leadsto Existence of a matching covering $\mathcal{B}_{2\epsilon}(W)$ is equivalent to (b).

$$\forall X' \subseteq_{\text{fin.}} \mathcal{B}_{2\epsilon}(W), \text{ we have } |X'| \leq |\cup_{x \in X'} N_G(x)|.$$

Stability theorem: Outline Step 1, 2, and 3

Let V and W be P -persistence modules.

V and W are ϵ -interleaved.

Cf. [Bje, '21, Ex. 5.3]
 \nRightarrow

- $\forall X' \subseteq \mathcal{B}_{2\epsilon}(V), |X'| \leq |\cup_{x \in X'} N_G(x)|$ holds.
- $\forall X' \subseteq_{\text{fin.}} \mathcal{B}_{2\epsilon}(W), |X'| \leq |\cup_{x \in X'} N_G(x)|$ holds.

Prop. (3)
 \iff
Hall

- \exists a matching in $G(\mathcal{B}_{2\epsilon}(V), \mathcal{B}(W))$ that covers $\mathcal{B}_{2\epsilon}(V)$.
- \exists a matching in $G(\mathcal{B}(V), \mathcal{B}_{2\epsilon}(W))$ that covers $\mathcal{B}_{2\epsilon}(W)$.

Prop. (2)
 \implies

\exists a matching in $G = (V, W; \epsilon)$ that covers $\mathcal{B}_{2\epsilon}(V) \sqcup \mathcal{B}_{2\epsilon}(W)$.

Prop. (1)
 \iff

V and W are ϵ -matched.

Stability theorem: Outline Step 4

Let V and W be **B -persistence modules**.

V and W are ϵ -interleaved

- B -persistence modules are interval-decomposable, with each interval determined by two elements of B .
- Λ_ϵ^B is a poset isomorphism ($\forall \epsilon \in \mathbb{R}_{\geq 0}$)

[T, '25]

\Rightarrow

- $\forall X' \subseteq \mathcal{B}_{2\epsilon}(V), |X'| \leq |\cup_{x \in X'} N_G(x)|$ holds.
- $\forall X' \subseteq_{\text{fin.}} \mathcal{B}_{2\epsilon}(W), |X'| \leq |\cup_{x \in X'} N_G(x)|$ holds.

Prop. (3)

\Leftrightarrow

Hall

- \exists a matching in $G(\mathcal{B}_{2\epsilon}(V), \mathcal{B}(W))$ that covers $\mathcal{B}_{2\epsilon}(V)$.
- \exists a matching in $G(\mathcal{B}(V), \mathcal{B}_{2\epsilon}(W))$ that covers $\mathcal{B}_{2\epsilon}(W)$.

Prop. (2)

\Rightarrow

\exists a matching in $G = (V, W; \epsilon)$ that covers $\mathcal{B}_{2\epsilon}(V) \sqcup \mathcal{B}_{2\epsilon}(W)$.

Prop. (1)

\Leftrightarrow

V and W are ϵ -matched.

$$\leadsto d_B^{\Lambda_\epsilon^B}(\mathcal{B}(V), \mathcal{B}(W)) = d_I^{\Lambda_\epsilon^B}(V, W)$$

Summary

- We introduce bipath PH, which is an extension of standard PH.
- Bipath PDs have stability (arXiv: 2503.01614), and it is shown using an isometry theorem.

Discussion

- Application of bipath PH to real data. \leadsto We recently discussed the use of it for image data analysis with material scientists.
 - Relation with interleaving distance for finite bipath posets by Alonso and Liu (arXiv: 2501.00322).
- \leadsto Universality of interleaving distance.

Thank you for your listening.