MathBio Workshop: Shape and Movement in Life Sciences

## A Computation of Bipath Persistent Homology and Bipath Persistence Diagrams

Shunsuke Tada Graduate School of Human Development and Environment 2025/03/26 Self-Introduction

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Research: Persistent homology using representation theory

Next Affiliation: Postdoc at Tohoku University (MathCCS)

Hobby: walking, running (Full-marathon experience), and others

An aim is talking about

# Bipath persistent homology

which is an extension of persistent homology.

- A Visualization/Computation of Bipath Persistence Diagrams (Joint work with Toshitaka Aoki, Emerson G. Escolar)
- Stability of bipath persistence diagrams (bipath PDs)

• T. Aoki, E. G. Escolar, and S. Tada. Bipath persistence. Japan Journal of Industrial and Applied Mathematics, 42:453–486, 2025.

• S. Tada. Stability of Bipath Persistence Diagrams. arXiv: 2503.01614, 2025.



(1)Introduction: Persistent Homology and others

(2) Bipath persistence diagrams/Computation

(3) Stability of bipath persistence diagrams

## Persistent homology

- *Persistent homology* (PH) is a tool in Topological Data Analysis.
- It captures the persistence of "shape" (connected components, holes or voids) of data by a *persistence diagram* (PD).





## Applications

- Material science
- Evolutional biology
- Cosmology (cosmic web)
- Computational gastronomy and others...

• Yasuaki Hiraoka Hiraoka, Takenobu Nakamura, Akihiko Hirata, Emerson G. Escolar, Kaname Matsue, and Yasumasa Nishiura. Hierarchical structures of amorphous solids characterized by persistent homology. *Proceedings of the National Academy of Sciences*, *113*(26), 7035-7040, 2016.

• Joseph Minhow Chan, Gunnar Carlsson, and Raul Rabadan. Topology of viral evolution. Proceedings of the National Academy of Sciences, 110(46):18566–18571, 2013

• Thierry Sousbie. The persistent cosmic web and its filamentary structure–I. theory and implementation. Monthly Notices of the Royal Astronomical Society, 414(1):350–383, 2011.

and others.

• Emerson G. Escolar, Yuta Shimada, and Masahiro Yuasa. A topological analysis of the space of recipes. International Journal of Gastronomy and Food Science, 39:101088, 2025.

### How to capture shape

PH: Persistent homology PD: Persistent diagram



a larger hole.

#### How to capture shape





Other settings: zigzag PH



$$S: S_1 \supset S_2 \subset S_3 \supset S_4 \subset S_5 \qquad H_i(S) = H_i(S_{r_1}) \leftarrow H_i(S_{r_2}) \rightarrow H_i(S_{r_3}) \leftarrow H_i(S_{r_4}) \rightarrow H_i(S_{r_5})$$
  
$$\cong \bigoplus I[b_i, d_i]^{m_{b_i, d_i}} \text{ (Interval-decomposable)}$$
  
Zigzag filtration

e.g., Temporal network



Gunnar Carlsson, and Vin De Silva. Zigzag persistence. Foundations of computational mathematics 10 (2010): 367-405. Myers, A., Muñoz, D., Khasawneh, F. A., & Munch, E. (2023). Temporal network analysis using zigzag persistence. *EPJ Data Science*, *12*(1), 6.

### Other settings: multiparameter PH



• density

• time

• partial charge, and others...

Gunnar Carlsson, and Afra Zomorodian. The Theory of Multidimensional Persistence. Discrete Comput Geom 42, 71–93 (2009).

### Difficulty of multiparameter PH

An indecomposable module: This is not interval.



Mickaël Buchet, and Emerson G. Escolar. Every 1D Persistence Module is a Restriction of Some Indecomposable 2D Persistence Module. Journal of Applied and Computational Topology.

### Difficulty of multiparameter PH

It is complicated to classify all the indecomposable module. (wild representation type)

1	1	1	1	1	1
$\cdot \rightarrow$	• →	• →	• →	$\cdot \rightarrow$	$\cdot \rightarrow$
1	1	1	1	1	1
• →	• →	• →	• →	• →	$\cdot \rightarrow$
1	1	1	1	1	1
$\cdot \rightarrow$	• →	$\cdot \rightarrow$	$\cdot \rightarrow$	$\cdot \rightarrow$	$\cdot \rightarrow$
1	1	1	1	1	1
$\cdot \rightarrow$	• →	• →	• →	$\cdot \rightarrow$	$\cdot \rightarrow$
1	1	1	1	1	1
• →	• →	• →	• →	• →	• →

Zbigniew Leszczyński. On the representation type of tensor product algebras. Fundamenta Mathematicae, 144(2):143–161, 1994.
 Zbigniew Leszczyński and Andrzej Skowroński. Tame triangular matrix algebras. Colloquium Mathematicum, 86(2):259 303, 2000.
 Ulrich Bauer, Magnus B. Botnan, Steffen Oppermann, and Johan Steen. Cotorsion torsion triples and the representation theory of filtered hierarchical clustering. Advances in Mathematics, 369:107171, 2020.





Do we have other arrangement of spaces like standard/zigzag filtration?



### **Bipath filtration** ([Aoki-Escolar-T, 25])



A pair of two filtrations sharing the same spaces at the ends.

We can consider a *bipath persistent homology* (bipath PH) of a filtration.



Bipath PH

We can get a *Bipath Persistence Diagram* (Bipath PD).





Multisets of intervals

PH: Persistent homology

->What can bipath PH do?

Bipath PH can be used to...

(1) study the persistence of topological features across a pair of filtrations connected at their ends, to compare the two filtrations.

$$\begin{array}{cccc} S_{\hat{0}} \hookrightarrow S_{1} \hookrightarrow S_{2} \hookrightarrow \cdots \hookrightarrow S_{n} \hookrightarrow S_{\hat{1}} \\ S : & \| & & \| \\ & S_{\hat{0}} \hookrightarrow S_{1'} \hookrightarrow S_{2'} \hookrightarrow \cdots \hookrightarrow S_{m'} \hookrightarrow S_{\hat{1}} \end{array}$$

Bipath PH can be used to...

(2) obtain an invariant for multiparameter PH.



Bipath PH can be used to...

(2) obtain an invariant for multiparameter PH.



Bipath filtration.

 $S_{1.4} \rightarrow S_{2.4} \rightarrow S_{3.4} \rightarrow S_{4.4}$  $S_{1,2} \rightarrow S_{2,2} \rightarrow S_{3,2} \rightarrow S_{4,2}$  $\uparrow \qquad \uparrow \qquad \uparrow \qquad \uparrow \qquad \uparrow$  $S_{1,1} \rightarrow S_{2,1} \rightarrow S_{3,1} \rightarrow S_{4,1}$ 

Bifiltration.

Bipath PH can be used to...

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Bipath filtration.

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Bifiltration.

Bipath PH can be used to...

(2) obtain an invariant for multiparameter PH.





		Bipath PH
	Interval-Decomposability	0
(1)	Visualization(Bipath PD)	0
(2)	Algorithm(Implementation)	0
(3)	Stability theorem for Bipath PD	0
	Application	_



		Bipath PH
	Interval-Decomposability	0
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## Bipath Persistence Diagram (Bipath PD)



How to visualize.

### Recall: standard persistence diagram



In the standard setting, intervals are visualized through the correspondence between intervals and points on the plane.  $[b,d) \mapsto (b,d) \in \mathbb{R}^2$ 

### **Bipath Persistence**

#### **Definition** Bipath poset

Let *m* and *n* be non-negative numbers. A *bipath poset*  $B_{n,m}$  is a poset consisting of two totally ordered sets

$$1 \leq 2 \leq \cdots \leq n$$
, and  $1' \leq 2' \leq \cdots \leq m'$ 

with the global minimum and the global maximum

 $\widehat{0} \text{ and } \widehat{1}.$ 

The Hasse diagram:  $B_{n,m}: \hat{0} \xrightarrow{1} \rightarrow 2 \xrightarrow{\leq} \cdots \rightarrow n$  $B_{n,m}: \hat{0} \xrightarrow{1} \rightarrow 2' \xrightarrow{\leq} \cdots \rightarrow n' \hat{1}$ 

### Bipath Persistence Intervals in $B \coloneqq B_{n,m}$ are one of the following forms:



• Each interval in B (except for B) is written by the pair (s, t)  $(s, t \in B)$  by taking end points of the interval.

### A correspondence of intervals and points

(1)Put elements of  $B_{n,m}$  (in clockwise) on the vertical and horizontal axes.



(2)Plot a point on the upper left region " $\mathcal{B}$ " for the interval  $B_{n.m}$ .

(3)Plot a point (s, t)for the interval  $\langle s, t \rangle$ .



Examples of bipath PD



• Compute bipath PD of  $H_1(S; k = \mathbb{F}_2)$ .



$$= \{ \langle 1, 1' \rangle, \langle 1, 2' \rangle, \langle 2, \hat{1} \rangle, \langle 3, 3 \rangle, \langle 1', 1' \rangle \}.$$
(A) (B) (C) (D) (E)



$$(A) (B) (C) (D) (E)$$



$$= \{ \langle 1, 1' \rangle, \langle 1, 2' \rangle, \langle 2, \hat{1} \rangle, \langle 3, 3 \rangle, \langle 1', 1' \rangle \}.$$
(A) (B) (C) (D) (E)











1st.

1st.



1st.



Bipath filtration of image data ( $30 \times 30$  pixel)



Colors are changed by the multiplicity of intervals.



Bipath filtration of image data ( $30 \times 30$  pixel)



Colors are changed by the multiplicity of intervals.

1st.

**〈1,2'**〉

1



Bipath filtration of image data ( $30 \times 30$  pixel)

Colors are changed by the multiplicity of intervals.

2

 $2^{\prime}$ 

11

1st.



Bipath filtration of image data ( $30 \times 30$  pixel)

î



Colors are changed by the multiplicity of intervals.

1st.



Bipath filtration of image data ( $30 \times 30$  pixel)

î



Colors are changed by the multiplicity of intervals.

1st.

1'

2'

 $\mathbf{2}$ 

1

Ő

î



Bipath filtration of image data ( $30 \times 30$  pixel)

Colors are changed by the multiplicity of intervals.

2

1

2'

1'

### Implementation: Computing Bipath PD



Implementation: Computing bipath PD

We gave a software for computing bipath PD on GitHub. (https://github.com/ShunsukeTada1357/Bipathposets) Input Output



#### Remark.

The computational algorithm is given in [Aoki-Escolar-T, Algorithm 2, 25].

### Implementation: Computing bipath PD

#### Note.

• We get a bipath PD by 2 times of standard algorithm for PH and matrix operations (whose size depend on the number of intervals).

$$\begin{array}{c|c} \text{Standard algorithm} \checkmark & \begin{array}{c} S_{\hat{0}} \hookrightarrow S_{1} \hookrightarrow S_{2} \hookrightarrow \cdots \hookrightarrow S_{n} \hookrightarrow S_{\hat{1}} \\ \parallel & \parallel \\ S_{\hat{0}} \hookrightarrow S_{1'} \hookrightarrow S_{2'} \hookrightarrow \cdots \hookrightarrow S_{m'} \hookrightarrow S_{\hat{1}} \end{array}$$

->Bipath PD can be computed without much more effort than the standard algorithm for PH.

### Stability of bipath persistence diagrams

### Background: Stability theorem for standard PH



Background: Stability theorem for standard PH **Stability theorem** (see [Frédéric Chazal, et al., 2009] for example) Let f and g be a real-valued function on top. sp. X. Then, we have  $d_{\mathrm{B}}(\mathcal{B}(V_f), \mathcal{B}(V_g)) \leq ||f - g||_{\infty}.$ 

- $\mathcal{B}(V_f)$  : The standard PD of PH by sublevelset filtration of f.
- $\mathcal{B}(V_g)$  : The standard PD of PH by sublevelset filtration of g.
- ·  $d_{\rm B}$ : Bottleneck distance between PDs.
- $||f g||_{\infty}$ : =  $\sup_{x \in X} |f(x) g(x)|$ .

#### It gives mathematical justification of the use of PH for studying noisy data.

David Cohen-Steiner, Herbert Edelsbrunner, and John Harer. Stability of persistence diagrams. Discrete& Computational Geometry, 37:103–120, 2007.
 Frédéric Chazal, David Cohen-Steiner, Marc Glisse, Leonidas J Guibas, and Steve Y Oudot. Proximityof persistence modules and their diagrams. In Proceedings of the twenty-fifth annual symposium onComputational geometry, pages 237–246, 2009.

### Stability theorem for bipath PD

**Theorem** [T, 25](Stability theorem for bipath PD). Let  $f = (f_1, f_2)$  and  $g = (g_1, g_2)$  be bipath functions satisfying ( $\blacklozenge$ ). Then, we have the following inequality:

 $d_{\mathrm{B}}(\mathcal{B}(V_f), \mathcal{B}(V_g)) \le ||f - g||_{\infty}.$ 



### Stability theorem for bipath PD: Example using implementation

- *f* : bipath function of a simplicial complex.
- f' = f + noise such that  $||f f'||_{\infty} = 1$  with ( $\blacklozenge$ ).



overlay PDs

### Stability theorem for bipath PD: Setting

We consider continuous version of a bipath poset B to discuss stability. (We have a natural embedding  $\iota: B_{n,m} \hookrightarrow B$ )



#### **Definition** (Bipath function)

A *bipath function* f on top. sp. X is a pair of B-valued functions  $f_1$  and  $f_2$  on X such that

$$\operatorname{Im} f_i \subseteq (\mathbb{R} \times \{i\}) \sqcup \{\pm \infty\} \subseteq B$$
  
and  
$$f_1^{-1}(-\infty) = f_2^{-1}(-\infty).$$

We denote by  $f: X \to B$  a bipath function on X.



#### **Definition** (Bipath sublevelset filtration)

Let f be a bipath function on a top. sp. X. For any b in B, let

$$(f \le b) := \begin{cases} X & \text{if } b = +\infty \\ f_1^{-1}(\{-\infty\}) & \text{if } b = -\infty \\ f_1^{-1}([-\infty, r] \times \{1\}) & \text{if } b = (r, 1) \\ f_2^{-1}([-\infty, r] \times \{2\}) & \text{if } b = (r, 2) \end{cases}$$

Then, they give a functor  $(f \leq \cdot) : B \rightarrow \text{Top.}$  We call it *bipath sublevelset filtration*.

$$(f \le +\infty) = X(=f_1^{-1}(\bigcirc)) \qquad \mathbb{R} \times \{1\}$$
$$-\infty \qquad \mathbb{R} \times \{2\} \qquad b +\infty$$

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$$(f \le -\infty) = f_1^{-1} (\bullet = \{-\infty\}) \underbrace{b}_{-\infty} \underbrace{\mathbb{R} \times \{1\}}_{\mathbb{R} \times \{2\}} \bullet +\infty$$

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Then, they give a functor  $(f \leq \cdot) : B \rightarrow \text{Top.}$  We call it *bipath sublevelset filtration*.

- $f: X \to B$ : A bipath function.
- $V_f := H_q(\cdot; k) \circ (f \leq \cdot)$ : Bipath PH of the sublevelset filtration of f.
- • $\mathcal{B}(V_f)$ :Bipath persistence diagram.

### Stability theorem for bipath PD: Theorem

- $f: X \rightarrow B$ : A bipath function.
- $V_f := H_q(\cdot; k) \circ (f \leq \cdot)$ : Bipath PH of the sublevelset filtration of f.

#### Condition

For two bipath functions f and g, we set the following condition:  $f_1^{-1}(\{-\infty\}) \stackrel{\text{def}}{=} f_2^{-1}(\{-\infty\}) \stackrel{\bigstar}{=} g_1^{-1}(\{-\infty\}) \stackrel{\text{def}}{=} g_2^{-1}(\{-\infty\}) \quad (\bigstar)$ 

**Theorem** [T, 25](Stability theorem for bipath PD). Let  $f = (f_1, f_2)$  and  $g = (g_1, g_2)$  be bipath functions satisfying ( $\blacklozenge$ ). Then, we have the following inequality:

$$d_{\mathrm{B}}(\mathcal{B}(V_f), \mathcal{B}(V_g)) \le ||f - g||_{\infty}.$$

•
$$||f - g||_{\infty} := \max\{||f_1 - g_1||_{\infty}, ||f_2 - g_2||_{\infty}\}.$$

### Summary We proposed a new setting called bipath PH, which is an extension of standard PH.

	Bipath PH
Interval-Decomposability	0
Visualization(Bipath PD)	0
Algorithm(Implementation)	0
Stability theorem for Bipath PD	Ο
Application	_

### Discussion

- Settings where bipath PH is essentially needed.
- Faster algorithms are needed.
- Is an assumption on stability theorem strong? Thank you for listening!

### Stability theorem for bipath PD: Bottleneck distance

#### Notation

•  $\Lambda := {\Lambda_{\epsilon} : B \to B}_{\epsilon \in \mathbb{R}}$ : A family of order-preserving maps satisfying:

$$\Lambda_{\epsilon}(\pm\infty) = \pm\infty$$
, and  $\Lambda_{\epsilon}((r,i)) = (r+\epsilon,i)$  for  $(r,i) \in \mathbb{R} \times \{i\}$ .



### Stability theorem for bipath PD: Bottleneck distance

#### **Proposition**

For an interval *I*, we have  $I \cap \Lambda_{-\epsilon}(I) = \emptyset$  if and only if

$$I \subseteq \mathbb{R} \times \{1\}$$
 with  $I \cap \{(r - \epsilon, 1) \mid (r, 1) \in I\} = \emptyset$ , or  $I \subseteq \mathbb{R} \times \{2\}$  with  $I \cap \{(r - \epsilon, 2) \mid (r, 2) \in I\} = \emptyset$ .

In this case, we say that I is  $\underline{\epsilon}$ -trivial.

$$\xrightarrow{\epsilon} I \xrightarrow{\epsilon} I$$

### Stability theorem for bipath PD: Bottleneck distance

#### **Definition** (*\epsilon*-matching)

Let X and Y be multisets of intervals. We define  $\epsilon$ -matching between X and Y by a partial matching  $\sigma: X \supseteq X' \xrightarrow{1:1} Y' \subseteq Y$  satisfying:

- Every  $I \in X \setminus X'$  is  $2\epsilon$ -trivial.
- Every  $J \in Y \setminus Y'$  is  $2\epsilon$ -trivial.
- If  $\sigma(I) = J$ , then  $I \subseteq \operatorname{Ex}^{\epsilon}(J)$  and  $J \subseteq \operatorname{Ex}^{\epsilon}(I)$ .

#### **Definition** (Bottleneck distance)

The bottleneck distance  $d_{\rm B}$  between a multisets of intervals in B is given by

 $d_{\mathcal{B}}(X,Y) := \inf\{\epsilon \ge 0 \mid \exists \epsilon \text{-matching between } X \text{ and } Y\}$ 

for multisets of intervals X and Y in B.

### **Theorem** [Aoki-Escolar-T, 23]

Let P be a connected finite poset. The following are equivalent.

(a) Every persistence module V over P is interval-decomposable.(b) The Hasse diagram of P is one of the following forms:

$$A_n(a): 1 \longleftrightarrow \cdots \longleftrightarrow n \qquad \qquad B_{n,m}: \begin{array}{c} 1 \to 2^{-} \to \cdots \to n \\ 0 & & \\ & & \\ & & \\ 1' \to 2' \to \cdots \to m' \end{array}$$

Zigzag posets (type A)

**Bipath posets**